

Computer Security

11. Firewalls & VPNs

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1

Conversation Isolation

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2

Fundamental Layer 2 & 3 Problems

- IP relies on store-and-forward networking
 - Network data passes through untrusted hosts
 - Routes may be altered to pass data through malicious hosts
- Packets can be sniffed
- TCP session state can be examined or guessed ...
... and TCP sessions can be hijacked
- No source authentication on IP packets

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Solution: Use private networks

Connect multiple geographically-separated private subnetworks together



But this is expensive ... and not feasible in many cases (e.g., cloud servers)

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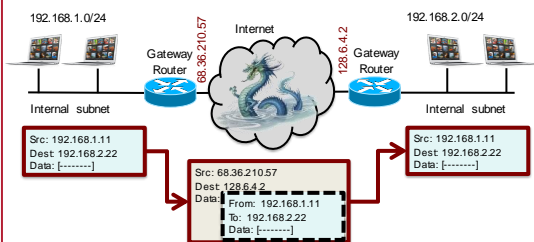
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Tunneling

Tunnel = Packet encapsulation

Treat an entire IP datagram as payload on the public network



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Tunnel mode vs. transport mode

- Tunnel mode
 - Communication between gateways: *network-to-network*
 - Or *host-to-network*
 - Entire datagram is encapsulated
- Transport mode
 - Communication between hosts
 - IP header is not modified

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IPsec

- Internet Protocol Security
- End-to-end solution at the IP layer
- Two protocols:
 - **IP Authentication Header Protocol (AH)**
 - Authentication & integrity of payload and header
 - **Encapsulating Security Payload (ESP)**
 - AH + Confidentiality of payload

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IPsec Authentication Header (AH)

Guarantees integrity & authenticity of IP packets

- MAC for the contents of the entire IP packet
- Over unchangeable IP datagram fields (e.g., not TTL or fragmentation)

Protects from:

- Tampering
- Forging addresses
- Replay attacks (signed sequence number in AH)

Layered directly on top of IP (protocol 51) - not UDP or TCP

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IPsec Encapsulating Security Payload (ESP)

Encrypts entire payload

- Plus authentication of payload + IP header (everything AH does) (may be optionally disabled - but you don't want to)

Directly on top of IP (protocol 51) - not UDP or TCP

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IPsec algorithms

- Integrity protection & authenticity
 - HMAC-SHA1
 - HMAC-SHA2
- Confidentiality
 - 3DES-CBC
 - AES-CBC
- Authentication
 - Kerberos, certificates, or pre-shared key authentication
- Key generation
 - Diffie-Hellman to exchange keying material for key generation
 - Key lifetimes determine when new keys are regenerated
 - Perfect forward secrecy
 - "Main mode master key PFS" - requires reauthentication & is CPU-intensive
 - "Quick mode session key PFS" - no reauthentication

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Conversation Isolation: Transport Layer SSL/TLS

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Transport Layer Security

- Provide a transport layer security protocol
- After setup, applications feel like they are using TCP sockets
 - SSL: Secure Socket Layer
- Created with HTTP in mind
 - Web sessions should be secure
 - Mutual authentication is usually not needed
 - Client needs to identify the server but the server won't know all clients
 - Rely on passwords after the secure channel is set up
- SSL evolved to TLS (Transport Layer Security)
 - SSL 3.0 was the last version of SSL ... and is considered insecure
 - We use TLS now ... but often still call it SSL

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TLS Protocol

- Goal
 - Provide authentication (usually one-way), privacy, & data integrity between two applications
- Principles
 - Use symmetric cryptography to encrypt data
 - Keys generated uniquely at the start of each session
 - Include a MAC with transmitted data to ensure message integrity
 - Use public key cryptography & X.509 certificates for authentication
 - Optional – can authenticate 0, 1, or both parties
 - Support many different key exchange, encryption, integrity, & authentication protocols – negotiate what to use at the start of a session

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TLS Protocol & Ciphers

Two sub-protocols

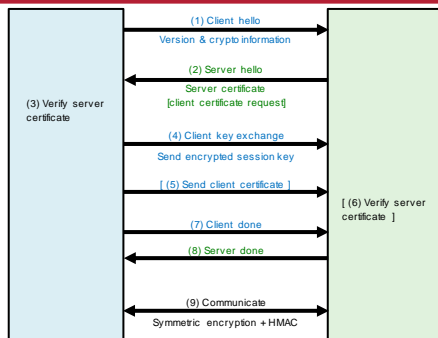
1. Authenticate & establish key
 2. Communicate
 - HMAC used for message authentication
- Key exchange
 - Public keys (RSA or Elliptic Curve)
 - Diffie Hellman keys
 - Ephemeral Diffie-Hellman keys (generated for each session)
 - Pre-shared key
- Data encryption
 - AES GCM, AES CBC, ARIA (GCM/CBC), ChaCha20-Poly1305, ...
- Data integrity
 - HMAC-MD5, HMAC-SHA1, HMAC-SHA256/384, ...

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TLS Protocol



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Benefits of TLS

- Benefits
 - Protects integrity of communications
 - Protects the privacy of communications
 - Validates the authenticity of the server (if you trust the CA)

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Problems with TLS

- Attacks
 - **Man-in-the-middle: BEAST attack in TLS 1.0**
 - Attacker was able to see Initialization Vector (IV) for CBC and deduce plaintext (known HTML headers & cookies)
 - Fixed by using explicit IVs for each new block
 - **Man-in-the-middle: crypto renegotiation**
 - Attacker can renegotiate the handshake protocol to disable encryption
 - Proposed fix: have client & server verify info about previous handshakes
 - **THC-SSL-DoS attack**
 - Attacker initiates a TLS handshake & requests a renegotiation of the encryption key – repeat over & over, using up server resources

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Problems with TLS

- Client authentication Problem
 - Client authentication is almost never used
 - Generating keys & obtaining certificates is not an easy process
 - Any site can request the certificate: user will be unaware anonymity is lost
 - Moving private keys around can be difficult (what about public systems?)
 - We usually rely on other authentication mechanisms (usually user name and password)

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Firewalls

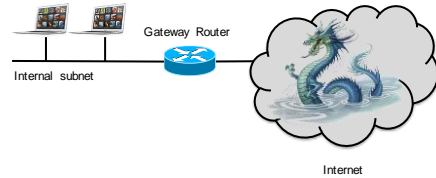
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Network Security Goals

- **Confidentiality**: sensitive data & systems not accessible
- **Integrity**: data not modified during transmission
- **Availability**: systems should remain accessible



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Firewall

- Separate your local network from the Internet
 - Protect the border between trusted internal networks and the untrusted Internet
- Approaches
 - Packet filters
 - Application proxies
 - Intrusion detection / intrusion protection systems

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Screening router

- **Border router** (gateway router)
 - Router between the internal network(s) and external network(s)
 - Any traffic between internal & external networks passes through the border router

Instead of just routing the packet, decide whether to route it

- **Screening router** = Packet filter
 - Allow or deny packets based on
 - Incoming interface, outgoing interface
 - Source IP address, destination IP address
 - Source TCP/UDP port, destination TCP/UDP port, ICMP command
 - Protocol (e.g., TCP, UDP, ICMP, IGMP, RSVP, etc.)

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Filter chaining

- An IP packet entering a router is matched against a set of rules: **access control list (ACL)** or **chain**
- Each rule contains criteria and an action
 - **Criteria**: packet screening rule
 - **Actions**
 - *Accept* – and stop processing additional rules
 - *Drop* – discard the packet and stop processing additional rules
 - *Reject* – and send an error to the sender (ICMP Destination Unreachable)
 - **Also**
 - *Route* – reroute packets
 - *Nat* – perform network address translation
 - *Log* – record the activity

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Filter structure is vendor specific

Examples

- Windows
 - *Allow*, *Block*
 - Options such as
 - Discard all traffic except packets allowed by filters (*default deny*)
 - Pass through all traffic except packets prohibited by filters (*default allow*)
- OpenBSD
 - *Pass* (allow), *Block*
- Linux *nftables* (*netfilter*)
 - Chain types: *filter*, *route*, *nat*
 - Chain control
 - *Return* – stop traversing a chain
 - *Jump* – jump to another chain (*goto* = same but no return)

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Network Ingress Filtering: incoming packets

Basic firewalling principle

Never have a direct inbound connection from the originating host from the Internet to an internal host – all traffic must flow through a firewall and be inspected

- Determine which services you want to expose to the Internet
 - e.g., HTTP & HTTPS: TCP ports 80 and 443
- Create a list of services and allow only those inbound ports and protocols to the machines hosting the services.
- **Default Deny** model - by default, "deny all"
 - Anything not specifically permitted is dropped
 - May want to log denies to identify who is attempting access

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Network Ingress Filtering

- Disallow IP source address spoofing
 - Restrict forged traffic (RFC 2827)
- At the ISP
 - Filter upstream traffic - prohibit an attacker from sending traffic from forged IP addresses
 - Attacker must use a valid, reachable source address
- Disallow incoming/outgoing traffic from private, non-routable IP addresses
 - Helps with **DDoS attacks** such as SYN flooding from lots of invalid addresses

```
access-list 199 deny ip 192.168.0.0 0.0.255.255 any log
access-list 199 deny ip 224.0.0.0 0.0.0.255 any log
.....
access-list 199 permit ip any any
```

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Network Egress Filtering (outbound)

- Usually we don't worry about outbound traffic.
 - Communication from a higher security network (internal) to a lower security network (Internet) is usually fine
- Why might we want to restrict it?
 - Consider: if a web server is compromised & all outbound traffic is allowed, it can connect to an external server and download more malicious code ... or launch a DoS attack on the internal network
 - Also, log which servers are trying to access external addresses

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Stateful Inspection

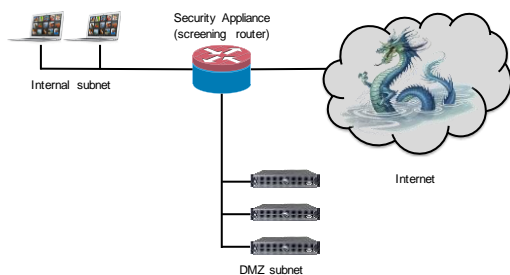
- Retain state information about a stream of related packets
- Examples
 - TCP connection tracking
 - Disallow TCP data packets unless a connection is set up
 - ICMP echo-reply
 - Allow ICMP echo-reply only if a corresponding echo request was sent
 - Related traffic
 - Identify & allow traffic that is related to a connection
 - Example: related ports in FTP

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Network Design: DMZ



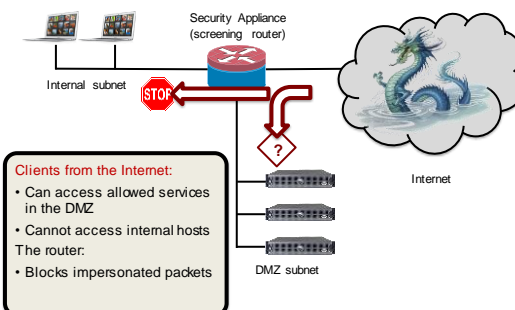
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Network Design: DMZ



Clients from the Internet:

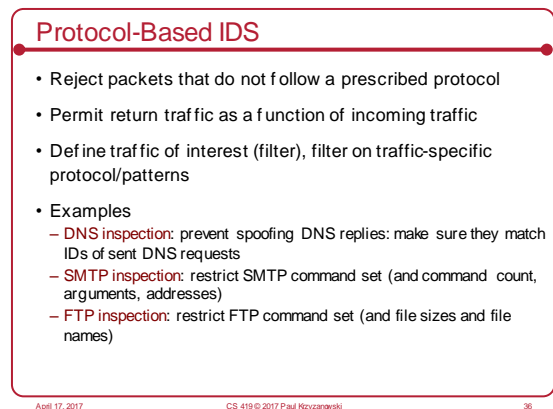
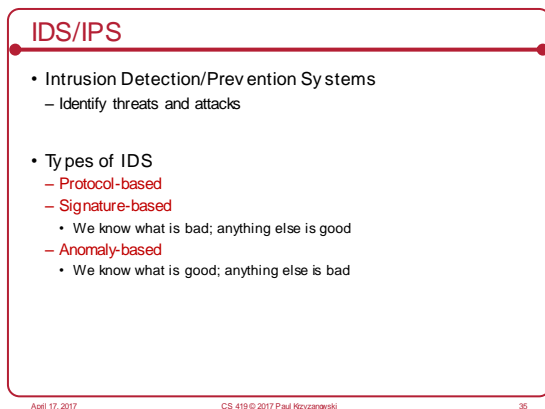
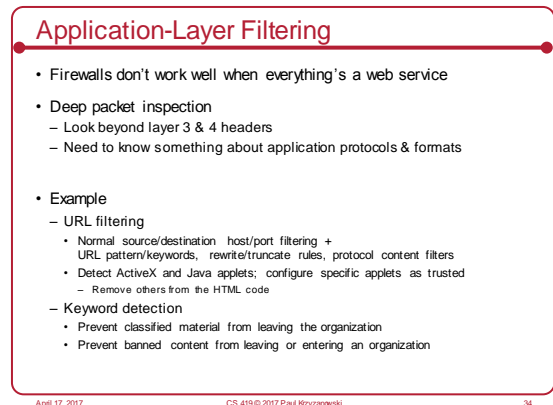
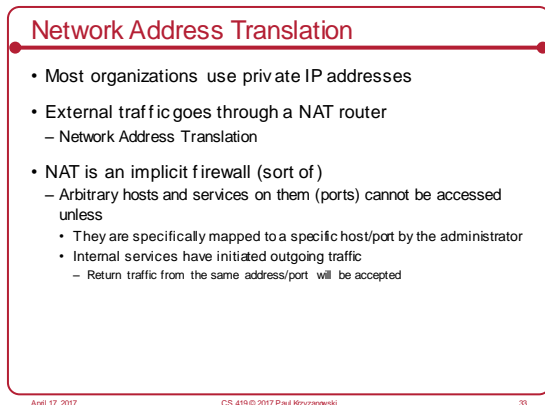
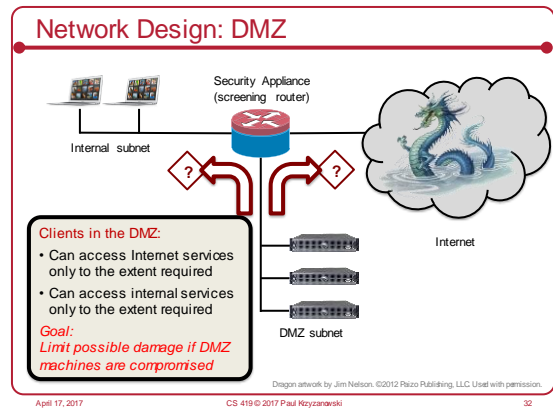
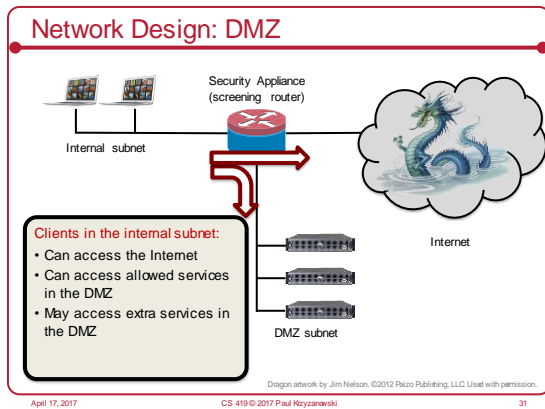
- Can access allowed services in the DMZ
- Cannot access internal hosts
- The router blocks impersonated packets

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Signature-based IDS

- Don't search for protocol violations but for exploits in programming
- Match patterns of known "bad" behavior
 - Viruses
 - Malformed URLs
 - Buffer overflow code

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Anomaly-based IDS

- Search for statistical deviations from normal behavior
 - Measure baseline behavior first
- Examples:
 - Port scanning
 - Imbalance in protocol distribution
 - Imbalance in service access

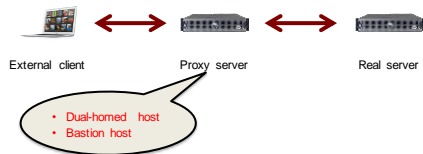
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Application proxies

- Proxy servers
 - Intermediaries between clients and servers
 - Stateful inspection and protocol validation



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Deperimeterization

- Boundaries & access between internal & external systems are harder to identify
 - Mobile systems
 - Cloud-based computing
 - USB flash memory
 - Web-based applications

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Host-based firewalls

- Run on the user's systems, not as dedicated firewalls
- Manage network-facing effects of malware
 - Allow only approved applications to send or receive data over the network
- Problem
 - If malware gets elevated privileges, it can reconfigure or disable the firewall
- Personal IDS
 - E.g., `fail2ban` on Linux
 - Scan log files to detect & ban suspicious IP addresses
 - High number of failed logins, probes, URLs that try to target exploits

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Intrusion detection & prevention problems

- There's a lot of stuff going on
 - People visit random websites with varying frequencies
 - Software accesses varying services
 - Buggy software may create bad packets
 - How do you detect what is hostile?
- Attack rates is miniscule ... compared to legitimate traffic
 - Even a small % of false positives can be annoying and hide true threats
- Environments are dynamic
 - Content from CDNs or other large server farms has a broad range of IP addresses
 - Malicious actors can coexist with legitimate ones

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Intrusion detection & prevention problems

- Encrypted traffic cannot be easily inspected
 - Just because you visit a web site using HTTPS doesn't mean the site is secure ... or hasn't been compromised
- Packet inspection is limiting
 - You may need to reconstruct sessions, which is time consuming
- Threats & services change
 - Rules have to be updated

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43

Summary

| | |
|------------------------------|--|
| Firewall (screening router) | 1 st generation packet filter that filters packets between networks. Blocks/accepts traffic based on IP addresses, ports, protocols |
| Stateful inspection firewall | Like a screening router but also takes into account TCP connection state and information from previous connections (e.g., related ports for TCP) |
| Application proxy | Gateway between two networks for a specific application. Prevents direct connections to the application from outside the network. Responsible for validating the protocol. |
| IDS/IPS | Can usually do what a stateful inspection firewall does + examine application-layer data for protocol attacks or malicious content |
| Host-based firewall | Typically screening router with per-application awareness. Sometimes includes anti-virus software for application-layer signature checking |
| Host-based IPS | Typically allows real-time blocking of remote hosts performing suspicious operations (port scanning, ssh logins) |

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DDoS

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DDoS: Distributed Denial of Service

- Compromise machines (create a botnet)
 - Use **amplification** techniques to generate a lot of traffic for targets
 - Exploit services that generate a lot of traffic to a small query
 - **DNS amplification:**
 - Small UDP query with forged source address results in large response
- Some targets were too huge to hurt with traffic
 - Amazon, Google, sites using CDNs such as Akamai
- Vast quantities of compromised systems reduce need for amplification
 - Create a botnet of millions of systems

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Dealing with DDoS

- Really difficult in general
- Bandwidth management routers
 - Either in data center or ISP
 - Limit outbound or inbound traffic on a per-IP basis
 - Detect DNS attack and set **null routing**
 - Traffic to attacked DNS goes nowhere
 - Egress filtering by ISPs
 - Attempt to find malicious hosts participating in DDoS or sending spam
 - Identify incoming attackers & block traffic at firewall
 - Difficult with a truly distributed DDoS attack

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The end

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